

WOLFRAM KRAUSE: OPTIMIZED NETWORK STRUCTURE AND ROUTING METRIC IN WIRELESS MULTIHOP AD HOC COMMUNICATION

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OUTLINE

WHAT IS BEING STUDIED?

Initial setting

Throughput as an optimization problem

CONTRIBUTION

Cumulative betweenness centrality

Optimization of transmission power

Optimization of the route

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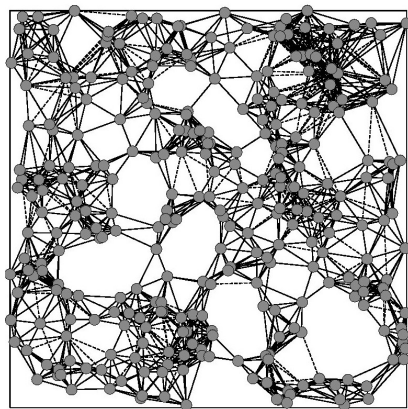
Optimization of the route

THE QUESTION

How to improve throughput of an wireless multihop ad hoc network?

THE MODEL

SOME ABSTRACTIONS ARE NEEDED FOR THE NETWORK



- unit square
- nodes homogeneously distributed

THE MODEL

WHAT MAKES THIS A WIRELESS NETWORK?

- undesigned network topology
- multihop -connections
- transmission power issues
- **agenda:** resolution model that can be applied in distributive fashion

THE MODEL

LIMITATIONS

- no mobility
- uniformly distributed traffic
- fairness is not an issue
- the throughput is

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WHAT IS END-TO-END THROUGHPUT?

- **information** transferred per timestep
- calculation
 - hard to do exactly
 - leads to simulations or approximations

HOW DOES THROUGHPUT SCALE?

- obeys scaling law $T_{e2e} = (N - N_0)^\gamma$
- theoretical maximum is $\gamma = 0.5$
- unoptimized network has $\gamma \approx 0.22$
- can be optimized to 0.41!
- why does it scale so badly?

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HOW TO OPTIMIZE THROUGHPUT?

- alter the transmission power
- alter the routes

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CUMULATIVE BETWEENNESS CENTRALITY

THE DEFINITION

WHAT IS BETWEENNESS CENTRALITY?

- a property of a node
- B_i number of end-to-end routes going through node i
 - The final node is not included to the path.
- note: depends on metric used in the network

HOW DOES IT CUMULATE?

- $B_i^{cum} = B_i + \sum_{j \in \mathcal{N}_i} B_j$, where \mathcal{N}_i is the neighbour nodes of i .
- cumulative sum of routes going through the node+neighbours of the node

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CBC, So WHAT?

WHAT IS THE USE OF THIS CONCEPT

- we can approximate: sending time = relative CBC
- why? the interference
- leads to $T_{e2e} = \frac{\text{nodepairs}}{\max_i \{B_i^{cum}\}}$
 - exact with fully connected network
 - exact with central hub

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TRANSMISSION POWER

BASIC PROBLEM WITH TRANSMISSION POWERS

WHAT TO OPTIMIZE?

- minimal power consumption vs. maximal connectivity
- maximal connectivity vs. minimal interference

- Short jumps and long routes?
- Short routes and long jumps?
- Leads to frustration.

TRANSMISSION POWER OPTIMIZATION

LET'S ITERATE A SOLUTION

INITIAL SETTING

- transmission power s.t. has atleast k_{min} neighbours
- A node can ask neighbours to speak louder, if needed
- power newer goes below initial setting
- \implies strong connectivity

THE OPTIMIZATION

- Gradient method search for local maxium
 - Test node by node for changes
- Small perturbations

TRANSMISSION POWER OPTIMIZATION

LET'S ITERATE A... BAD SOLUTION

- The results are poor:
 - first local maximum is the best
 - gets worse in every iteration
- **Even worse:** the structural properties make solutions useless!
 - effects are indirect
 - effects are unpredictable

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THINKING OUT OF THE BOX

FINDING A NEW DIRECTION

The previous approach failed – something has to be done

CUMULATIVE BETWEENNESS CENTRALITY

CBC IS USEFUL FOR ROUTING METRIC

- forget transmission powers
- don't minimize hops, avoid traffic
- leads to better utilization of the network

CBC ROUTING METRIC IN PRACTICE

HEY, THE CBC DEPENDS ON THE ROUTES ITSELF!

THE SETUP HAS TO BE DONE WITH CARE

- Initialize $B_k^{cum} = 1$
- Iterate
 - Node by node
 - dijkstra-like fashion
- Experimental result: two rounds are enough
- Leads to good results

SUMMARY

- The **cumulative betweenness centrality** is good for
 - making throughput estimates
 - an alternative routing metric
- Optimization by modifying of transmission powers is hard, because effects are indirect and delocalized.

FOR FURTHER READING I



W. Krause, J. Scholz, M. Greiner.

Optimized network structure and routing metric in wireless multihop ad hoc communication

[arXiv:cs.NI/0503010](https://arxiv.org/abs/cs.NI/0503010)



P. Gupta, P. R. Kumar,

The capacity of wireless networks

IEEE Trans. Info. Theory IT46 (2) (2000) 388404.